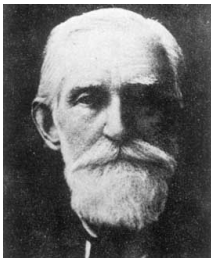


Chebyshev Approximation by Polynomials, Rational Functions and Splines



Pafnuti Lwowitsch Chebyshev (1821 - 1894)

Best Approximation

Best Approximation and Minimal deviation

Let $(F, \|\cdot\|)$ be a normed space and V a non-empty subset of F .
For given $f \in F$, determine $v^* \in V$, such that

$$\|f - v^*\| \leq \|f - v\| \text{ for all } v \in V.$$

v^* is the **best approximation** of f w.r.t. V

$\rho_V(f) = \|f - v^*\| (= \inf_{v \in V} \|f - v\|)$: **Minimal deviation**

Main Questions

Main Questions and Problems (depending on F , $\|\cdot\|$ and V):

1. Existence
2. Uniqueness
3. Characterization
4. (Numerical) construction, algorithms
5. Lower and upper bounds for $\varrho_V(f)$

In this talk: $F = C[a, b]$

In the first (and main) part: V is a Haar subspace¹ of $C[a, b]$.²

Definition

An m -dimensional subspace V of $C[a, b]$ is called **Haar subspace**, if each $v \in V$, $v \neq 0$, has at most $m - 1$ zeros in $[a, b]$.

¹Alfred Haar, 1885 - 1933

²Sometimes also Chebychev subspace

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Important examples

- Π_n , the space of polynomials of degree at most n , is a $(n + 1)$ -dimensional Haar subspace on each interval $[a, b]$.
- T_n , the space of trigonometric polynomials of degree at most n , is a $(2n + 1)$ -dimensional Haar subspace on $[0, 2\pi)$.

¹Alfred Haar, 1885 - 1933

²Sometimes also Chebychev subspace



L_1 -Approximation

Let, for all $h \in C[a, b]$,

$$\|h\|_1 = \int_a^b |h(x)| dx$$

denote the L_1 -**Norm**.

Theorem (Jackson 1930)

Let $f \in C[a, b]$ and let V be a Haar subspace of $C[a, b]$. Then f possesses one and only one best L_1 -approximation w.r.t. V .

Theorem (Kripke and Rivlin 1965)

Let $f \in C[a, b]$ and let V be a Haar subspace of $C[a, b]$. Then $v^* \in V$ is the best L_1 -approximation of f , if and only if

$$\int_a^b v(x) \cdot \operatorname{sgn}(f - v^*)(x) dx \leq \int_Z |v(x)| dx, \quad (1)$$

for all $v \in V$, where Z denotes the set of zeros of $(f - v^*)$. If, in addition, $(f - v^*)$ has only finitely many zeros, then (1) simply reads

$$\int_a^b v(x) \cdot \operatorname{sgn}(f - v^*)(x) dx = 0.$$



Chebyshev Approximation

Chebyshev norm

The norm

$$\|h\|_{\infty} = \max_{x \in [a, b]} |h(x)|$$

is called

- maximum norm
- uniform norm
- L_{∞} -norm
- Chebyshev norm

Chebyshev norm

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- Chebyshev norm

From now on throughout the talk!

Negative example (for uniqueness)

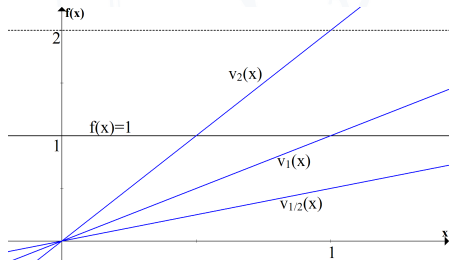
Let $[a, b] = [0, 1]$, $V = \text{span}\{x\}$ and $f(x) \equiv 1$.

Since $v(0) = 0$ for all $v \in V$, we have $\rho_V(f) \geq 1$.

Set $v_\alpha(x) = \alpha x$ with $0 \leq \alpha \leq 2$. Then

$$\|f - v_\alpha\| = \max_{x \in [0,1]} |(1 - \alpha x)| = 1.$$

So v_α is best approximation for all $0 \leq \alpha \leq 2$.



Theorem of Haar

Haar subspaces are the „right“ spaces for Chebyshev approximation:³

Theorem (Haar/Kolmogorov 1918/48)

Let V be an n -dimensional subspace of $C[a, b]$.

1. If V is a Haar subspace, then each $f \in C[a, b]$ possesses one and only one best approximation w.r.t. V .
2. If there exists a function $v \in V$, $v \neq 0$, with at least n zeros in $[a, b]$, then there exists an $\bar{f} \in C[a, b]$, which has more than one best approximation w.r.t. V .

³Andrej Nikolajevich Kolmogorov, 1903 - 1987

Characterization

Let $h \in C[a, b]$, and let $E(h)$ denote the set of **extremal points** of h :

$$E(h) = \{x \in [a, b]; |h(x)| = \|h\|_\infty\}.$$

Definition (Alternant)

Let $h \in C[a, b]$ and $m \in \mathbb{N}$. A set of points

$$a \leq x_1 < x_2 < \dots < x_m \leq b$$

is called an **alternant** of h of length m , if

$$\{x_1, x_2, \dots, x_m\} \subset E(h)$$

and

$$h(x_i) = -h(x_{i+1}) \text{ for } i = 1, 2, \dots, m - 1.$$

Characterization

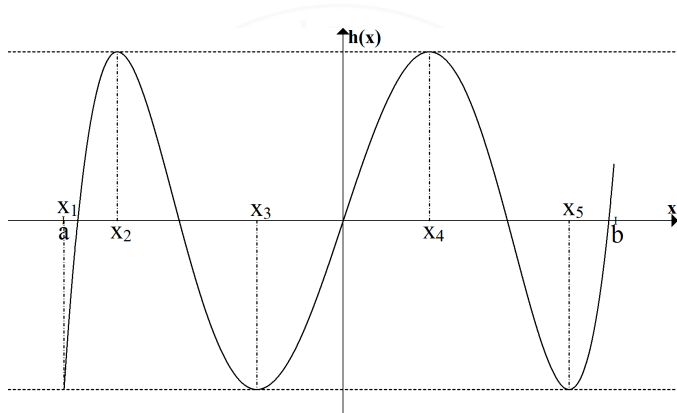


Figure: Alternant of length 5

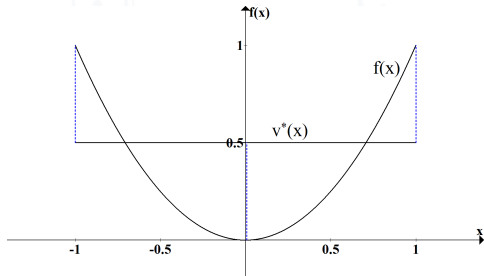
Characterization

Fundamental Theorem (Chebyshev 1859)

Let V be an n -dimensional Haar subspace of $C[a, b]$ and $f \in C[a, b]$.

Then $v^* \in V$ is the best approximation of f w.r.t. V if and only if the **error function** $(f - v^*)$ has an alternant of length $n + 1$.

Example: $f(x) = x^2$ on $[-1, 1]$, Π_1 ($\mapsto n = 2, n + 1 = 3$)



Characterization

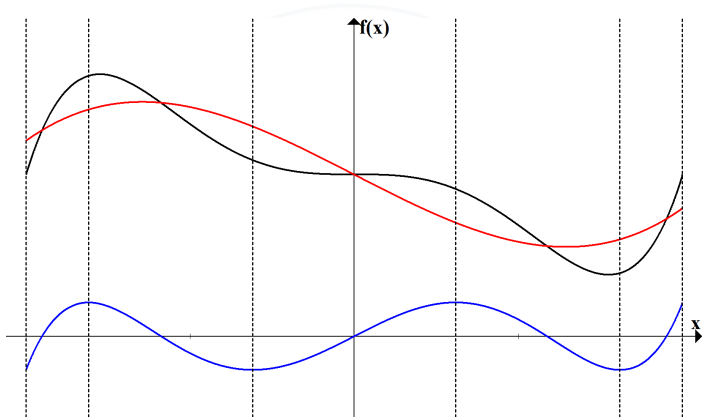


Figure: Black: Function $f(x)$. Red: Best Approximation $v^*(x)$. Blue: Error function $(f - v^*)(x)$

Chebyshev's Theorem is fundamental for the whole theory. Among others, it leads to

- estimates (inclusion) of the minimal deviation (de la Vallée Poussin⁴)
- iterative algorithm for the construction of the best approximation (Remez⁵)

⁴Charles Jean Gustave Nicolas Baron de la Vallée Poussin, 1866 – 1962

⁵Evgenii Yakovlevich Remez, 1895 – 1975

De la Vallée Poussin

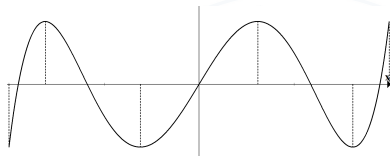


Figure: Alternant of error function

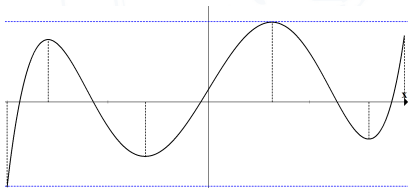


Figure: Almost Alternant of error function

Theorem (de la Vallée Poussin)

Let V be an n -dimensional Haar subspace of $C[a, b]$ and $f \in C[a, b]$. Assume that there exists a function $\tilde{v} \in V$ and $n + 1$ points

$$a \leq t_1 < t_2 < \dots < t_{n+1} \leq b$$

such that

$$(f - \tilde{v})(t_i) \neq 0 \text{ for } i = 1, 2, \dots, n + 1$$

and

$$\operatorname{sgn}\{(f - \tilde{v})(t_i)\} = -\operatorname{sgn}\{(f - \tilde{v})(t_{i+1})\} \text{ for } i = 1, 2, \dots, n.$$

Then

$$\min_{1 \leq i \leq n+1} |(f - \tilde{v})(t_i)| \leq \rho_V(f) \leq \|(f - \tilde{v})\|.$$

De la Vallée Poussin

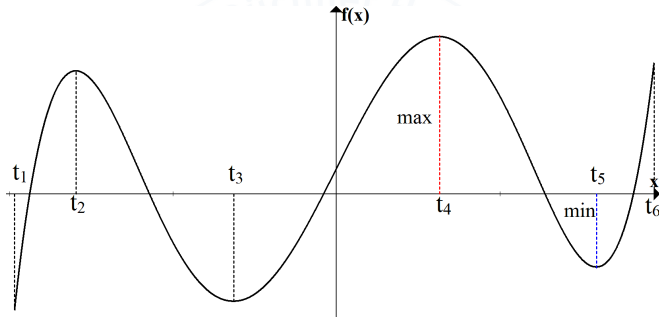


Figure: Illustration of de la Vallée Poussin

De la Vallée Poussin

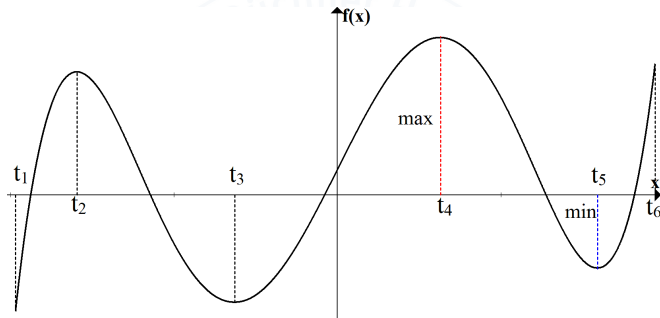


Figure: Illustration of de la Vallée Poussin

Leads to the fundamental **Remez algorithm**.

Remez algorithm

Basic idea: Iterative construction of

- a sequence of functions $v^k(x)$ with $\lim_{k \rightarrow \infty} v^k(x) = v^*(x)$, and simultaneously
- a sequence of numbers λ^k with $\lim_{k \rightarrow \infty} |\lambda^k| = \varrho_V(f)$

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- a sequence of functions $v^k(x)$ with $\lim_{k \rightarrow \infty} v^k(x) = v^*(x)$, and simultaneously
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Fundamental lemma

Let V be an n -dimensional Haar subspace of $C[a, b]$ and $f \in C[a, b]$. Then, for an arbitrary point set (a so-called **reference**)

$$R = \{a \leq x_1 < x_2 < \cdots < x_{n+1} \leq b\}$$

there exists a uniquely determined function $v \in V$ and a uniquely determined real number λ , such that

$$(f - v)(x_i) = (-1)^i \cdot \lambda \quad \text{for } i = 1, 2, \dots, n + 1.$$

Remez algorithm

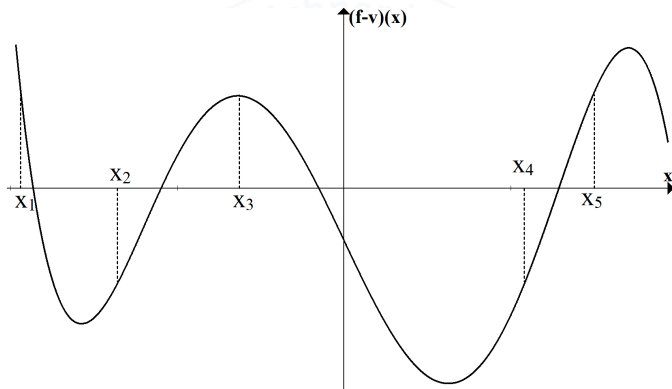


Figure: Illustration of the fundamental lemma: $n = 4$, $\lambda =$ length of dotted lines

Remez algorithm

Remez algorithm

- Step 0: Choose an arbitrary reference R^0 on $[a, b]$. Set $k = 0$.

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Remez algorithm

- Step 0: Choose an arbitrary reference R^0 on $[a, b]$. Set $k = 0$.
- Step 1: Determine $v^k(x)$ and λ^k according to the fundamental lemma.
 - If $\|f - v^k\| = |\lambda^k|$, then $v^k(x)$ is the best approximation and R^k is the alternant. Stop algorithm and praise the Lord.
 - If $\|f - v^k\| > |\lambda^k|$, but $(\|f - v^k\| - |\lambda^k|) < \varepsilon$, stop. Otherwise proceed with Step 2.

Remez algorithm

Remez algorithm

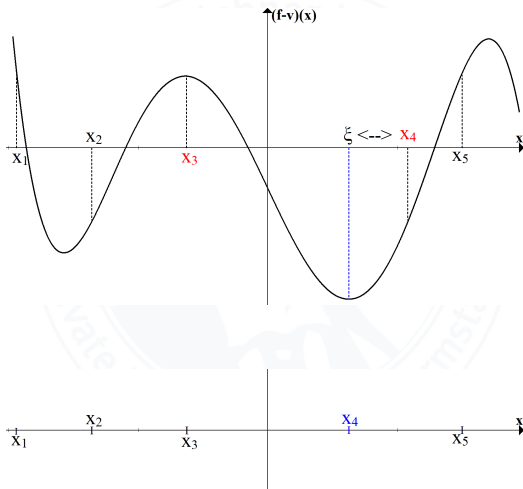
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- Step 2: Determine an extremal point ξ^k of $(f - v^k)$. Assume $x_i^k < \xi^k < x_{i+1}^k$ (cyclic). Replace x_i^k or x_{i+1}^k by ξ^k s.t. the alternating behaviour of the error function remains valid. This defines a new reference R^{k+1}

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- Step 3: Set $k := k + 1$ and proceed with Step 1.

Remez algorithm



Remez algorithm

Main Theorem (Remez)

With the notations and definitions from above, the following statements hold:

1. For all $k \in \mathbb{N}_0$:

$$|\lambda^k| < |\lambda^{k+1}| \leq \varrho_V(f)$$

2. $\lim_{k \rightarrow \infty} |\lambda^k| = \varrho_V(f)$
3. $\lim_{k \rightarrow \infty} \|v^k - v^*\| = 0.$

Remez algorithm: Numerical Example

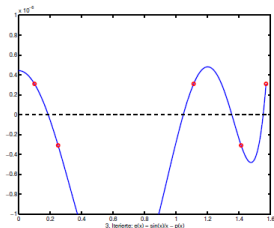
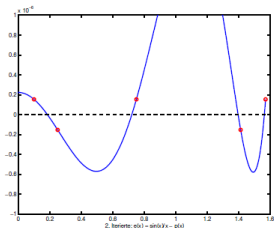
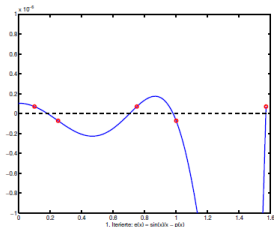
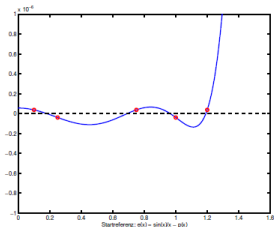
Example (due to H.J.Oberle): Approximate

$$f(x) = \frac{\sin(x)}{x} \quad \text{on} \quad \left[0, \frac{\pi}{2}\right]$$

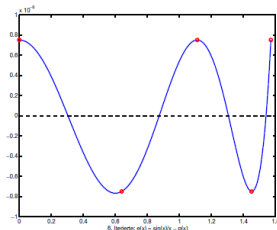
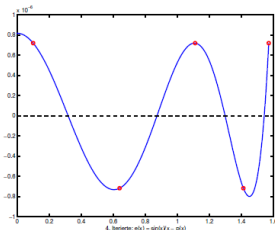
by even polynomials of degree 6, i.e.

$$p(x) = a_0 + a_1x^2 + a_2x^4 + a_3x^6.$$

Remez algorithm: Numerical Example



Remez algorithm: Numerical Example



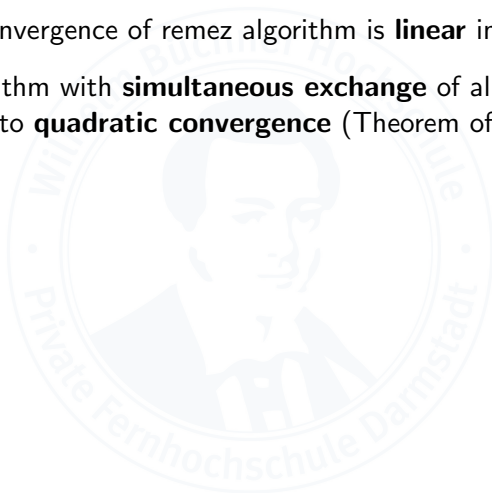
$$p^*(x) = 0.999999 - 0.166667x^2 + 0.008313x^4 - 0.000185x^6$$

$$\varrho_V(f) = 0.75439 \cdot 10^{-6}$$

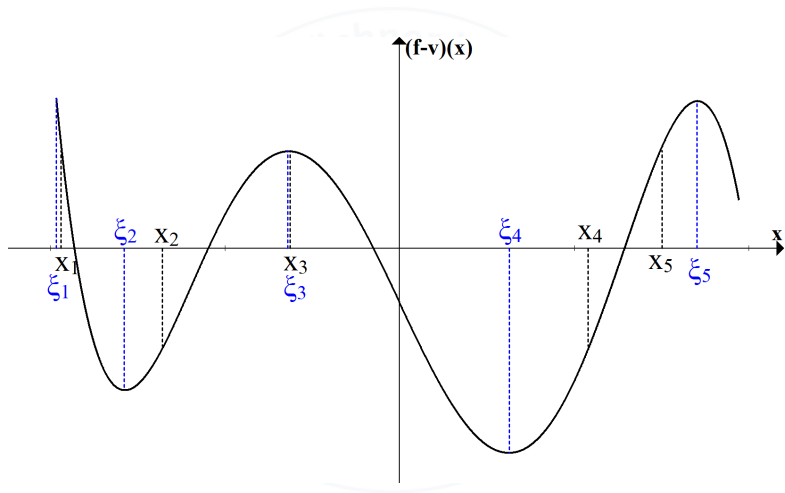
Remez algorithm with simultaneous exchange

Remark: Convergence of remez algorithm is **linear** in general.

Remez algorithm with **simultaneous exchange** of all reference points leads to **quadratic convergence** (Theorem of Veidinger 1960).



Remez algorithm with simultaneous exchange





Weak Haar Spaces

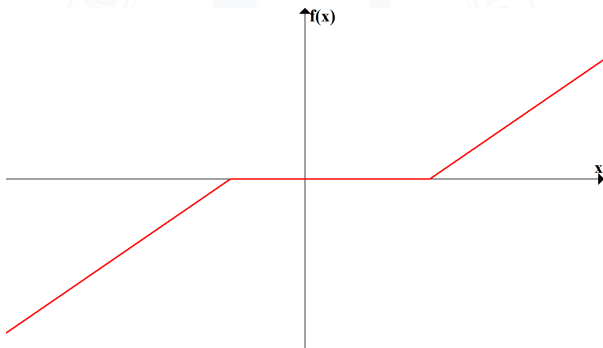
Definition (Weak Haar Space)

An m -dimensional subspace V of $C[a, b]$ is called **Weak Haar subspace**, if each $v \in V$ has at most $m - 1$ sign changes in $[a, b]$.

Weak Haar Space

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Theorem (Jones and Karlowitz 1970)

Let V be an n -dimensional Weak Haar subspace of $C[a, b]$ and $f \in C[a, b]$.

1. Then there exists a best approximation $v^* \in V$ of f w.r.t. V , such that the error function $(f - v^*)$ has an alternant of length $n + 1$.
2. If, for some $v^* \in V$, the error function $(f - v^*)$ has an alternant of length $n + 1$, then v^* is a best approximation of f .

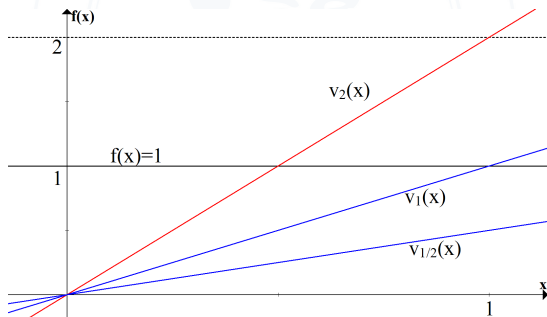
Remark: Uniqueness is lost!

Weak Haar Space

Example: Let $[a, b] = [0, 1]$, $V = \text{span}\{x\}$ and $f(x) \equiv 1$.

Set $v_\alpha(x) = \alpha x$ with $0 \leq \alpha \leq 2$. Then

- v_α is best approximation for all $0 \leq \alpha \leq 2$.
- $(f - v_2)$ has an alternant of length 2.



Definition (Spline)

With $k, m \in \mathbb{N}, m \geq 2$, let be given a set of $k + 1$ fixed **knots**

$$a = t_0 < t_1 < \dots < t_{k-1} < t_k = b$$

on $[a, b]$. A function s is called **spline of order m** (w.r.t. these knots), if

1. $s \in C^{m-2}[a, b]$
2. the restriction of s to each interval $[t_i, t_{i+1}]$ is a polynomial of degree $m - 1$.

Splines

Example: $m = 2$

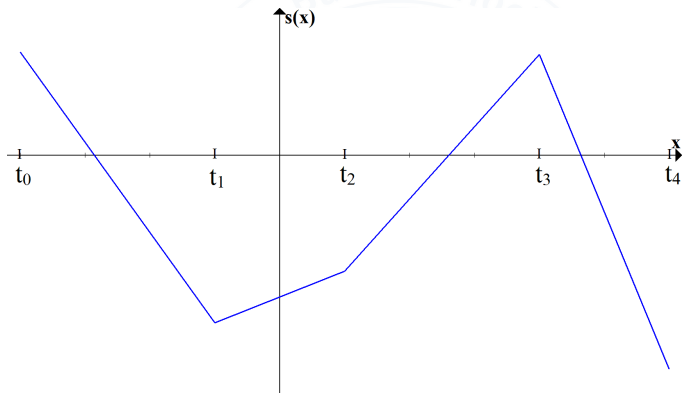


Figure: Linear spline

Splines

Example: $m = 3$

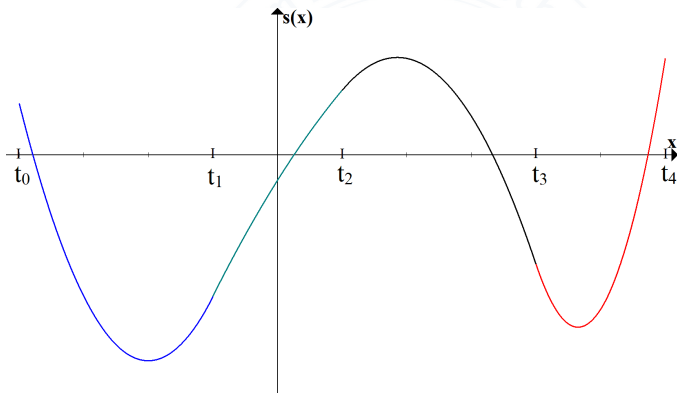


Figure: Quadratic spline

Theorem

Let

$$S_m(t_0, \dots, t_k)$$

denote the set of all splines of order m w.r.t. the above defined knot set.

Then $S_m(t_0, \dots, t_k)$ is a Weak Haar subspace of $C[a, b]$ of dimension

$$\dim S_m(t_0, \dots, t_k) = m + k - 1.$$

Splines: Characterization

For best spline approximations, there is a characterization theorem by Rice and Schumaker.⁶

Theorem (Rice and Schumaker 1968)

Let $f \in C[a, b]$. Then a spline $s^* \in S_m(t_0, \dots, t_k)$ is a best approximation of f , if and only if there is a knot interval $[t_p, t_{p+q}]$, such that the error function $(f - s^*)(x)$ has an alternant of length $m + q$ on $[t_p, t_{p+q}]$.

⁶John R. Rice, born 1934; Larry L. Schumaker, born 1939

Splines: Strong uniqueness

Uniqueness is very complicated; the following result is due to Nürnberger.⁷

Theorem (Nürnberger 1978)

Let $f \in C[a, b]$. Then a spline $s^* \in S_m(t_0, \dots, t_k)$ is the (strongly) unique best approximation of f , if and only if the following statements on the error function $(f - s^*)(x)$ are true:

⁷Günther Nürnberger, 1948 - 2013

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1. $(f - s^*)(x)$ has an alternant of length $m + k$ on $[a, b]$.
2. For each (relevant) j and ν , $(f - s^*)(x)$ has an alternant of length $j + 1$ on the intervals

$$[t_0, t_j), (t_{n-j}, t_n], (t_\nu, t_{\nu+m+j-1}).$$

⁷Günther Nürnberger, 1948 - 2013

Splines: Algorithm

There is also an **algorithm** for computing best spline approximations, due to Sommer⁸ and Nürnberger (1983).

Exchange Algorithm of **Remez type**, but much more complicated.

Two fundamental differences:

⁸Manfred Sommer, born 1950

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1. System of equations for computing new reference might be unsolvable → modified exchange rule

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Splines: Algorithm

There is also an **algorithm** for computing best spline approximations, due to Sommer⁸ and Nürnberger (1983).

Exchange Algorithm of **Remez type**, but much more complicated.

Two fundamental differences:

1. System of equations for computing new reference might be unsolvable → modified exchange rule
2. Spline approximations are not unique in general → convergence (proof) very complex

⁸Manfred Sommer, born 1950



Nonlinear Approximation

Nonlinear Case

So far: Linear approximation

Now: Nonlinear approximation (e.g.: rational functions, free knot splines)

Definition (Tangent Space)

Let $F \subset C[a, b]$ be a set of functions of the form $f(p, x)$, depending on a parameter vector $p = (p_1, p_2, \dots, p_k)$. The **tangent space** of F at the point $f(p, x)$ is the linear space

$$T(p) = \text{span} \left\{ \frac{\partial}{\partial p_1} f(p, x), \frac{\partial}{\partial p_2} f(p, x), \dots, \frac{\partial}{\partial p_k} f(p, x) \right\}$$

Set $d(p) = \dim(T(p))$.

Nonlinear Case

Definition (Haar condition)

The set F satisfies the

- **local Haar condition**, if for every parameter vector p the tangent space $T(p)$ is a Haar subspace

Nonlinear Case

Definition (Haar condition)

The set F satisfies the

- **local Haar condition**, if for every parameter vector p the tangent space $T(p)$ is a Haar subspace
- **global Haar condition**, if for every choice of parameter vectors p and q , $p \neq q$, the difference function

$$f(p, x) - f(q, x)$$

has at most $d(p) - 1$ zeros on $[a, b]$.

Nonlinear Case

Semi-uniqueness and characterization.⁹

Theorem (Meinardus 1961)

Let $f \in C[a, b]$. Suppose that F satisfies the local and the global Haar condition. Then the following holds:

1. f has at most one best approximation w.r.t. F .

⁹Günter Meinardus, 1926 - 2007

Nonlinear Case

Semi-uniqueness and characterization.⁹

Theorem (Meinardus 1961)

Let $f \in C[a, b]$. Suppose that F satisfies the local and the global Haar condition. Then the following holds:

1. f has at most one best approximation w.r.t. F .
2. The function $v^* \in F$ is the best approximation of f , if and only if the error function $(f - v^*)(x)$ has an alternant of length $d(p) + 1$.

⁹Günter Meinardus, 1926 - 2007

Rational Approximation

Special case: **Rational functions**

$$f(p, x) = \frac{P_m(x)}{Q_n(x)} = \frac{a_0 + a_1x + \dots + a_mx^m}{1 + b_1x + \dots + b_nx^n}$$

i.e.

$$p = (a_0, a_1, \dots, a_m, b_1, b_2, \dots, b_n)$$

Tangent Space:

$$\text{span} \left\{ \frac{1}{Q_n(x)}, \frac{x}{Q_n(x)}, \dots, \frac{x^m}{Q_n(x)}, \frac{-xP_m(x)}{Q_n^2(x)}, \dots, \frac{-x^n P_m(x)}{Q_n^2(x)} \right\}$$

Satisfies the local and global Haar condition!

Rational Approximation

Uniqueness and characterization:

Theorem (Chebyshev)

Let $f \in C[a, b]$ and $R_{m,n}$ the set of rational functions as above. Then the following holds:

1. f has one and only one best approximation w.r.t. $R_{m,n}$.

Rational Approximation

Uniqueness and characterization:

Theorem (Chebyshev)

Let $f \in C[a, b]$ and $R_{m,n}$ the set of rational functions as above. Then the following holds:

1. f has one and only one best approximation w.r.t. $R_{m,n}$.
2. The function $r^* \in R_{m,n}$ is the best approximation of f , if and only if the error function $(f - r^*)(x)$ has an alternant of length $l + 1$, where

$$l = \begin{cases} m + n + 1 - \min(m - m_1, n - n_1), & \text{if } P_m(x) \not\equiv 0 \\ m + 1, & \text{if } P_m(x) \equiv 0 \end{cases}$$

and m_1, n_1 denotes the exact degrees of P_m and Q_n resp.



Free Knot Splines

Free Knot Splines

So far: Splines with fixed knots. **Now:** Free knots.

Definition (Free Knot Splines)

A function $s : [a, b] \mapsto \mathbb{R}$ is called **spline of order m with $k - 1$ free knots**, if there exists a set of knots

$$a = t_0 \leq t_1 \leq \dots \leq t_{k-1} \leq t_k = b$$

s.t. $s \in S_m(t_0, \dots, t_k)$.

The set $S_{m,k}$ of all these splines is (highly) nonlinear, since the knots are parameters!

Free Knot Splines

The following result on the tangent space goes in principle back to Wulbert:¹⁰

Theorem (Wulbert 1973)

The tangent space of the spline set $S_{m,k}$ at the point s is the space of splines with the same knots, but (explicitly known) higher multiplicities.

So the tangent space is a weak Haar space!

Idea: Approximate it by Haar spaces.

¹⁰Daniel Wulbert, born 1940

Definition (Gauß transform)

For $h \in C(\mathbb{R})$ and $t > 0$, the mapping resp. the function

$$G(h, t)(x) = \frac{1}{\sqrt{4\pi t}} \int_{-\infty}^{\infty} h(u) \cdot \exp\left(-\frac{(u-x)^2}{4t}\right) du$$

is called **Gauß transform** of h with parameter t (if it exists).

Theorem

1. For all $x \in \mathbb{R}$:

$$\lim_{t \rightarrow 0} G(h, t)(x) = h(x).$$

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2. If

$$\text{span}\{h_1(x), h_2(x), \dots, h_m(x)\}$$

is a weak Haar subspace, then

$$\text{span}\{G(h_1, t)(x), G(h_2, t)(x), \dots, G(h_m, t)(x)\}$$

is a Haar subspace.

Free Knot Splines

We again will make use of the tangent space.
Important: Switch between Gauß transform and linearization.

Theorem (W. 1998)

Let $F \subset C[a, b]$ be a set of functions of the form $f(p, x)$, depending **continuously** on a parameter vector p .

Then

$$T(G(f(p, x))) = G(T(f(p, x)))$$

where as above T denotes the tangent space.

Free Knot Splines

Algorithm for free knot splines (Outline)

- Problem: Given $f \in C[a, b]$, determine best approximation s^* of f w.r.t. $S_{m,n}$.

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Characterization and uniqueness theorems, first numerical results: Meinardus and W. (2001)

Free Knot Splines

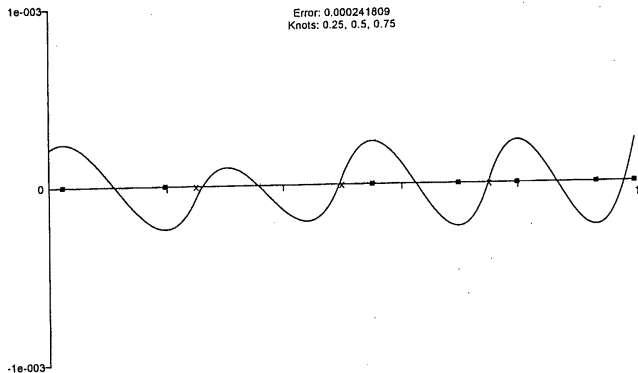


Figure: Approximation of $\exp(x)$: Fixed knots, $m = 3$, $k = 4$

Free Knot Splines

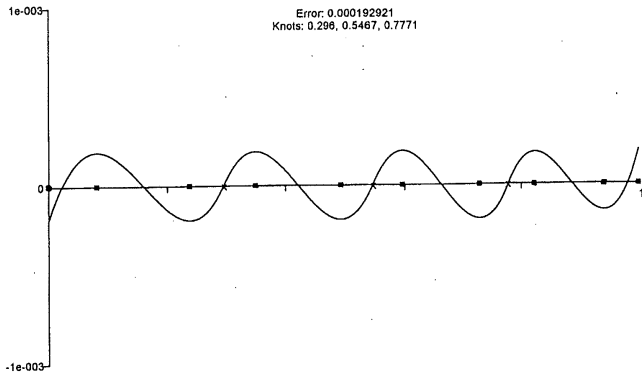


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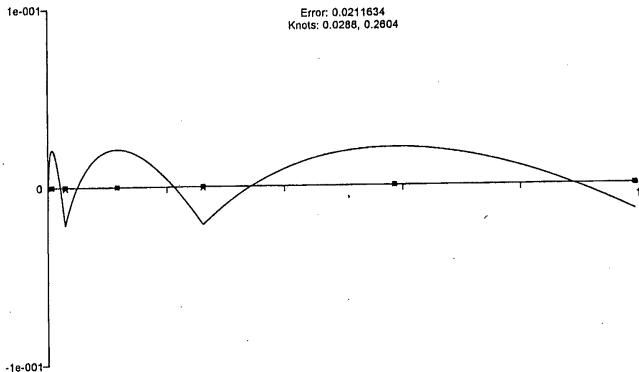


Figure: Approximation of \sqrt{x} : Free knots, $m = 2$, $k = 3$



Thank you very much for your attention

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Definition (Norm)

Given a vector space F . A **norm** $\| \cdot \|$ is a mapping from F into the real numbers with the following properties:

1. $\|f\| = 0 \implies f = 0$
2. $\|a \cdot f\| = |a| \cdot \|f\|$ for all $f \in F$ and $a \in \mathbb{R}$
3. $\|f_1 + f_2\| \leq \|f_1\| + \|f_2\|$ for all $f_1, f_2 \in F$

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Properties

1. $\|f\| \geq 0$ for all $f \in F$
2. $\|f\| = 0 \iff f = 0$

Strong uniqueness

Definition (Strong uniqueness)

Let V be a non-empty subset of $C[a, b]$ and $f \in C[a, b]$.
 $v^* \in V$ is called **strongly unique best approximation** of f , if there exists a constant $\gamma_f > 0$, such that

$$\|f - v\| \geq \|f - v^*\| + \gamma_f \|v - v^*\|$$

for all $v \in V$.

Remark: Implies Uniqueness!



L_2 -Approximation

Let

$$\|h\|_2 = \left(\int_a^b (h(x))^2 dx \right)^{\frac{1}{2}}$$

denote the L_2 -**Norm**, and set

$$L_2[a, b] = \left\{ h : [a, b] \mapsto \mathbb{R}; \int_a^b (h(x))^2 dx < \infty \right\}.$$

This norm is induced by an **inner product**: Define, for $g_1, g_2 \in C[a, b]$,

$$\langle g_1, g_2 \rangle = \int_a^b g_1(x) \cdot g_2(x) dx$$

Then

$$\|h\|_2 = \langle h, h \rangle^{\frac{1}{2}}$$

Theorem (Existence and Uniqueness)

Let V be a finite-dimensional subspace of $L_2[a, b]$ and $f \in L_2[a, b]$. Then f possesses one and only one best L_2 -approximation w.r.t. V .

Theorem (Characterization)

Let V be a finite-dimensional subspace of $L_2[a, b]$ and $f \in L_2[a, b]$. $v^* \in V$ is the best approximation of f w.r.t. V , if and only if

$$\langle f - v^*; v \rangle = 0 \text{ for all } v \in V.$$

v^* is the **orthogonal projection** of f onto V .

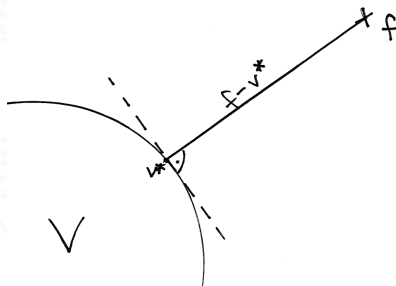


Figure: Orthogonal projection